

# Motion-Prediction-based Wireless Scheduling for Multi-User Panoramic Video Streaming

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# Motivation

- Emerging Commercial Head-mounted Displays (HMDs)
- Panoramic video streaming provides an immersive experience for users as if they are in a virtual 3D world

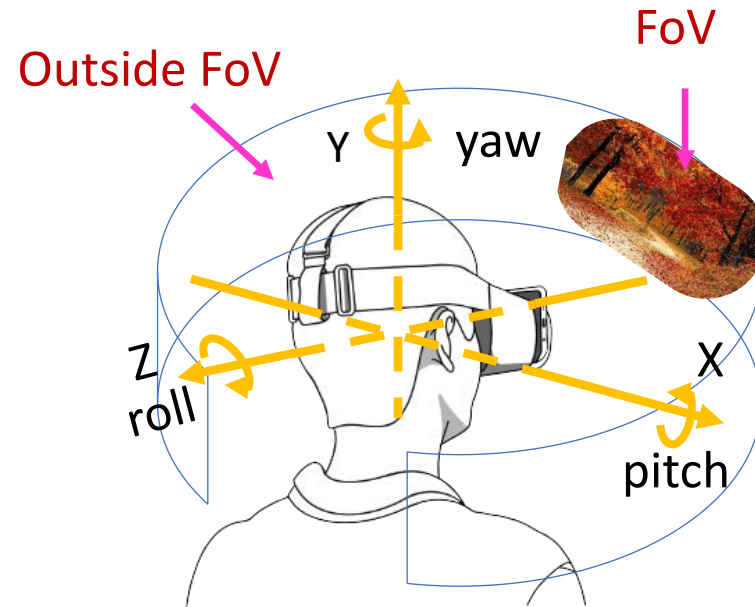


Virtual Reality Classroom

- **Main challenges:**
  - **Large network bandwidth requirement:** 4~6x bandwidth consumption of a regular video with the same resolution
  - **Seamless user experience:** users would compete for limited bandwidth

# Opportunity

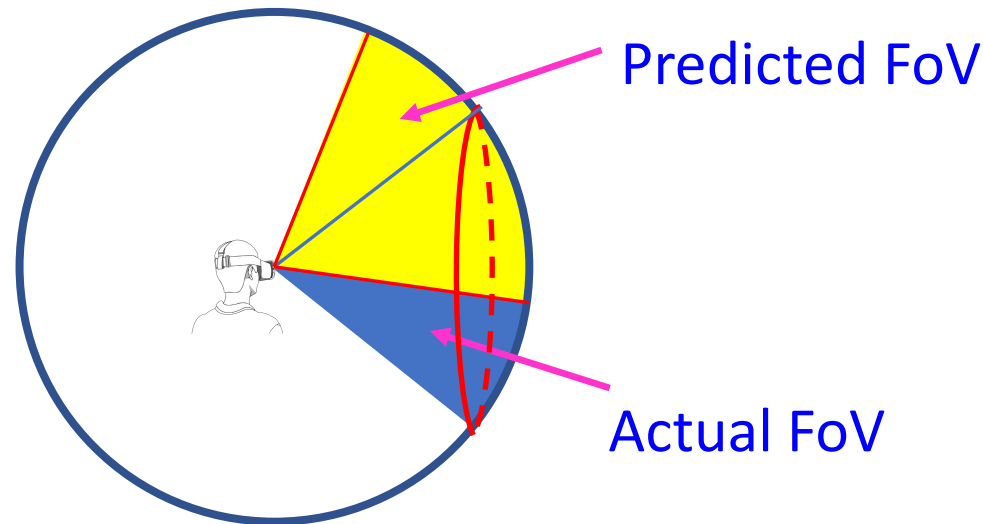
- A user may only see as little as 20% of 360° scenes, known as **Field of View (FoV)**. It is sufficient to deliver 20% of 360° video scenes under **perfect motion prediction**.



Rotation coordinates and FoV

# Practical Challenges

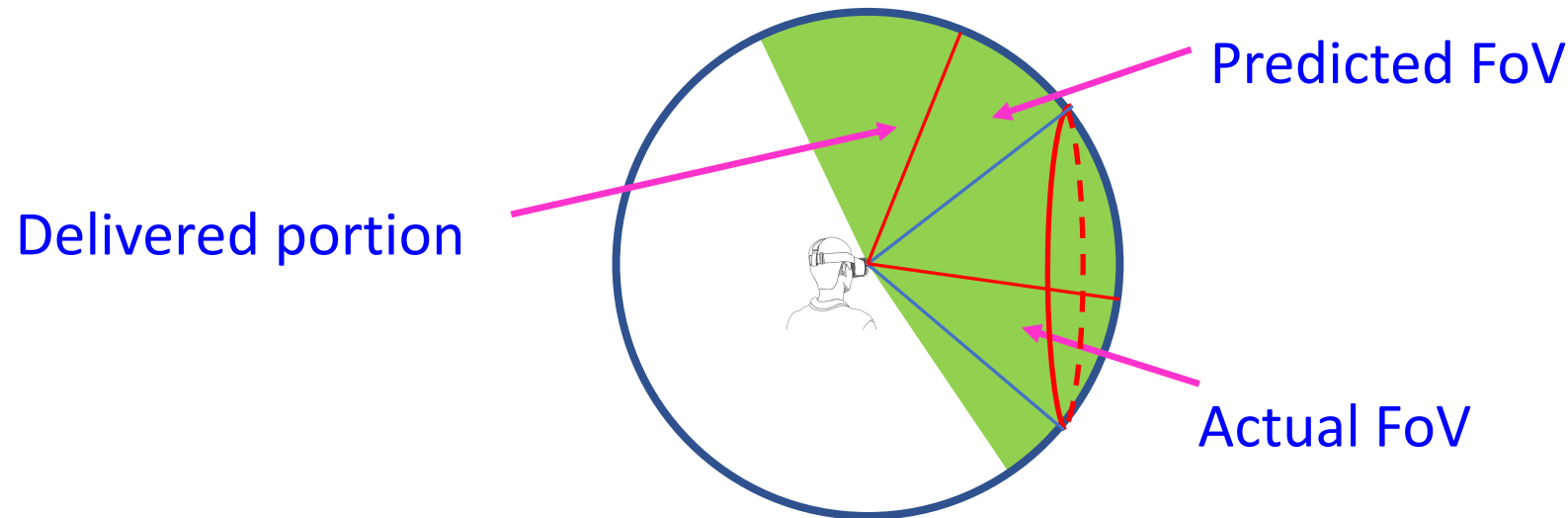
- Imperfect prediction: should deliver a portion larger than the FoV



With imperfect prediction, some of the actual FoV will be missed if we only send the predicted FoV

# Practical Challenges

- Imperfect prediction: should deliver a portion larger than the FoV



The actual FoV can be completely covered if we send a larger portion based on the predicted FoV

# Successful Viewing Probability

- Prediction errors of both pitch and yaw angles of user  $n$  follow the **normal distribution** with standard deviation  $\sigma_n^X$  and  $\sigma_n^Y$ , respectively
- We characterize the successful viewing probability  $\delta_n(S_n[t])$  as a function of the delivered portion ratio  $S_n[t]$  (normalized allocated rate)

$$\delta_n(S_n[t]) = \text{erf}^2 \left( \frac{\gamma_n(S_n[t])}{\sqrt{2}} \right)$$

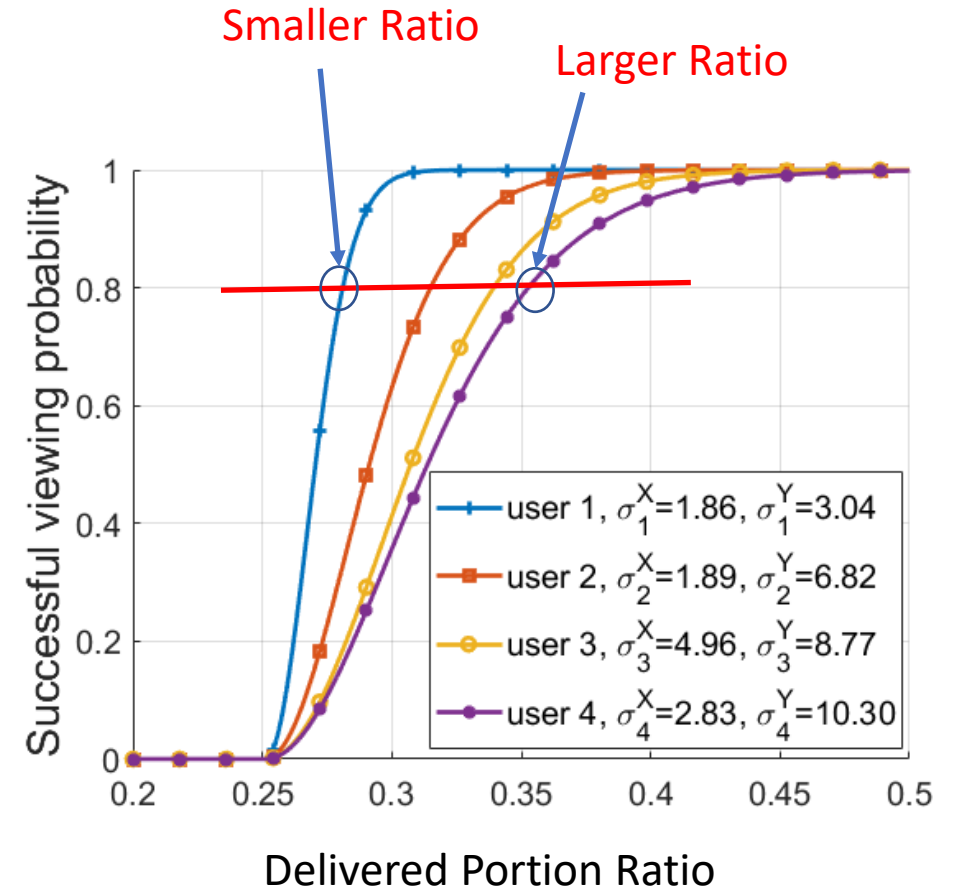
where  $\text{erf}(x) \triangleq \frac{2}{\pi} \int_0^x e^{-y^2} dy$ ,  $\gamma_n(S_n[t])$  is the number of standard deviations of the prediction error,

i.e.,

$$\begin{matrix} \hat{X}_n[t] - \gamma_n(S_n[t])\sigma_n^X < X_n[t] < \hat{X}_n[t] + \gamma_n(S_n[t])\sigma_n^X \\ \hat{Y}_n[t] - \gamma_n(S_n[t])\sigma_n^Y < Y_n[t] < \hat{Y}_n[t] + \gamma_n(S_n[t])\sigma_n^Y \end{matrix}$$

Predicted angles

Real angles



# First Goal: Maximizing Throughput

- Maximize the application-level throughput (defined as the weighted sum of the expected successful viewing probability)

$$\max_{(S_n[t])_{n=1}^N} \lim_{L \rightarrow \infty} \frac{1}{L} \sum_{t=0}^{L-1} \sum_{n=1}^N w_n E[\delta_n(S_n[t])]$$

- **Constraints:**

- Wireless interference constraints

$$\text{s.t. } (S_n[t])_{n=1}^N \in \mathbf{S}^{(C[t])}, \forall t \geq 0$$

- The average allocated transmission rate should not be less than some minimum rate

$$\lim_{L \rightarrow \infty} \frac{1}{L} \sum_{t=0}^{L-1} E[S_n[t]] \geq r_n, \forall n$$

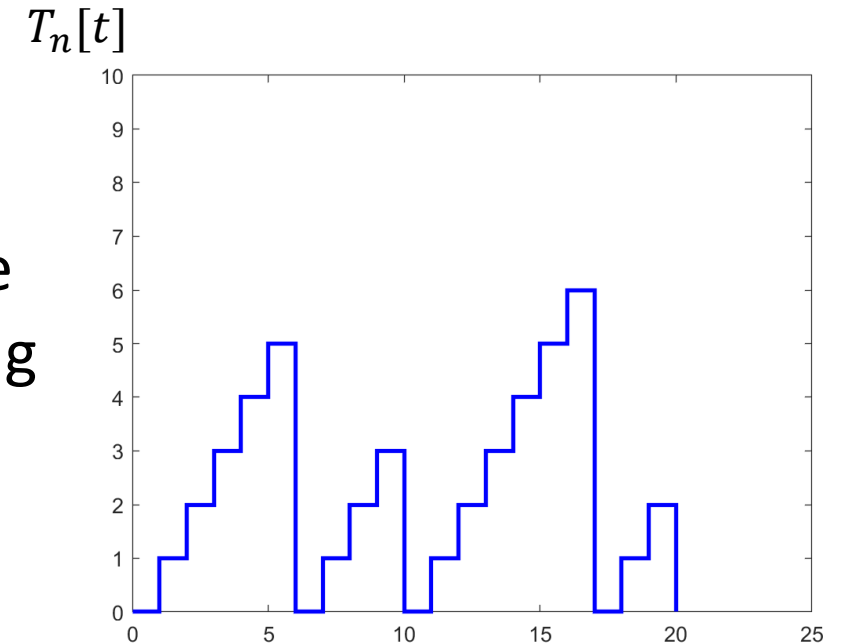
# Second Goal: Providing Seamless Experience

- Seamless user experience, keep service regularity (defined as the variance of the time between two consecutive successful views for each user)
- **Time-Since-Last-Service (TSLs) counter:**

$$T_n[t + 1] \triangleq \begin{cases} 0, & \text{if } I_n(S_n[t]) = 1; \\ T_n[t] + 1, & \text{otherwise.} \end{cases}$$

- [Li, Li, Eryilmaz 2014] showed that minimizing the expected TSLs counter is equivalent to minimizing the normalized variance of the time duration between successful services
- Our second goal is equivalent to minimizing

$$\lim_{L \rightarrow \infty} \frac{1}{L} \sum_{t=0}^{L-1} E[T_n[t]]$$



t



# Motivating Example

time rate user	0	1	2	3	...
User 1	1,	0,	0,	0,	...
User 2	0,	1,	0,	0,	...
User 3	0,	0,	1,	0,	...
User 4	0,	0,	0,	1,	...

(a) Service rate of each user in each time slot.

time result user	0	1	2	3	...
User 1	1,	0,	0,	0,	...
User 2	0,	1,	0,	0,	...
User 3	0,	0,	1,	0,	...
User 4	0,	0,	0,	1,	...

(b) Successful content delivery rate of each user in each time slot.

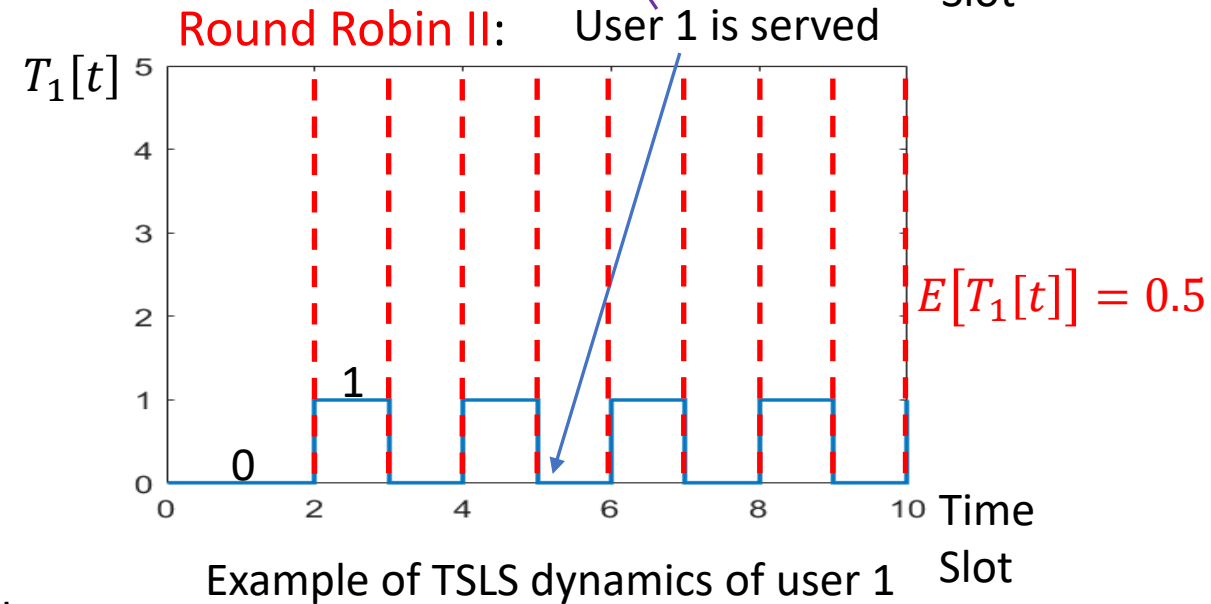
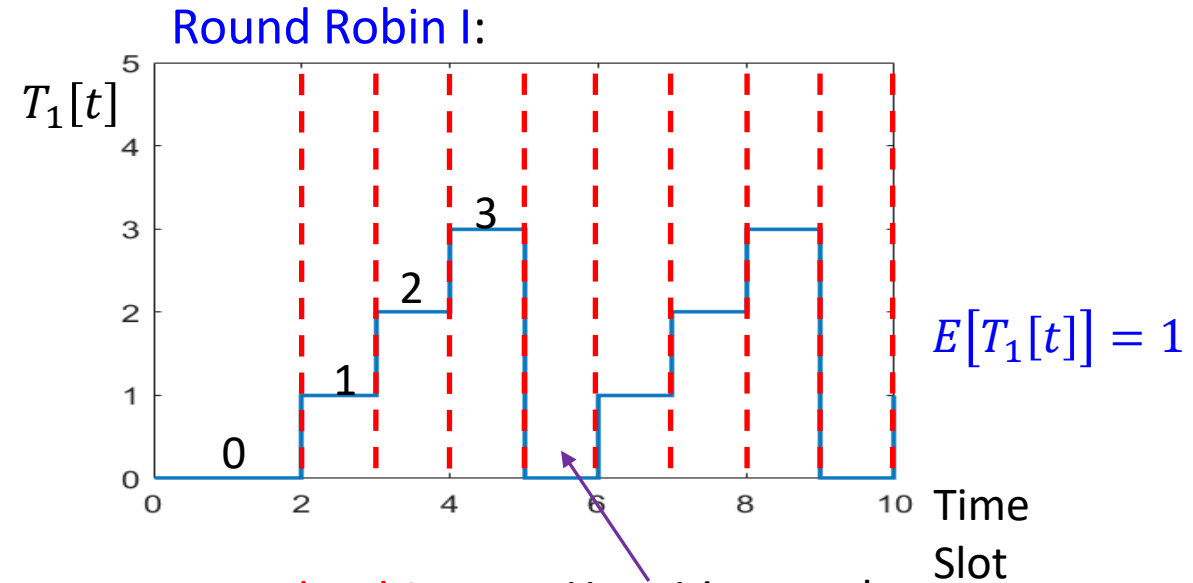
# Motivating Example

user \ time	rate	0	1	2	3	...	
User 1	1, 0.5	0, 0	0, 0.5	0, 0	0, 0.5	0, 0	...
User 2	0, 0.5	1, 0	0, 0.5	0, 0	0, 0.5	0, 0	...
User 3	0, 0	0, 0.5	1, 0	0, 0.5	0, 0	0, 0.5	...
User 4	0, 0	0, 0.5	0, 0	1, 0	0, 0.5	0, 0	...

(a) Service rate of each user in each time slot.

user \ result	time	0	1	2	3	...	
User 1	1, 1	0, 0	0, 1	0, 0	0, 1	0, 0	...
User 2	0, 1	1, 0	0, 0	0, 1	0, 0	0, 0	...
User 3	0, 0	0, 1	1, 0	0, 0	0, 1	0, 0	...
User 4	0, 0	0, 1	0, 0	1, 0	0, 0	1, 1	...

(b) Successful content delivery rate of each user in each time slot.



# Motion Prediction

- Perform independently for each user in each axis since the correlation between  $X_n[t]$  and  $Y_n[t]$  is much smaller than their autocorrelations

- Autoregressive Model:

$$\hat{X}_n[t] = - \sum_{k=1}^W a_n[k] X_n[t-k] \quad \text{and} \quad \hat{Y}_n[t] = - \sum_{k=1}^W b_n[k] Y_n[t-k]$$

- [Fuller 2009] showed that the prediction error converges to the **Gaussian distribution** as the number of data samples goes to infinity

# Scheduling Algorithm Design

- A virtual queue for each user that measures the degree of violation of the average service rate constraint

$$Q_n[t + 1] \triangleq (Q_n[t] + r_n - S_n[t])^+, \forall n, \forall t$$

- Non-standard Lyapunov function that combines the virtual queue and TSLs counter

$$V[t] = \frac{1}{2} \sum_{n=1}^N Q_n^2[t] + \eta \sum_{n=1}^N T_n[t]$$

- **Wireless scheduling:**

- Select the schedule  $\mathcal{S}^*[t]$  following:

$$\mathcal{S}^*[t] \in \operatorname{argmax}_{\mathcal{S} \in \mathcal{S}^{(c)}} \sum_{n=1}^N (S_n[t] Q_n[t] + (\eta T_n[t] + K w_n) \delta_n(S_n[t]))$$

where  $\eta$  and  $K$  are tunable parameters

Weight of User  $n$

Virtual Queue

TSLs counter

Successful Viewing Prob.

- Computation complexity: like the Max-Weight algorithm

# Operation Example

$$S^*[t] \in \operatorname{argmax}_{S \in \mathcal{S}^{(c)}} \sum_{n=1}^N \overbrace{(S_n[t]Q_n[t] + (\eta T_n[t] + K w_n) \delta_n(S_n[t]))}^{W_n[t]}$$

$$Q_1[1] = 0.5, T_1[1] = 3, w_1 = 0.3$$



User 1

$$Q_2[1] = 0.5, T_2[1] = 2, w_2 = 0.6$$



User 2

$$Q_3[1] = 0.5, T_3[1] = 1, w_3 = 1$$



User 3

Tunable parameters:  $\eta = 0.2, K = 1$

Schedule one user at each time slot

Available rate:  $S_n[t] = 0.5$

Successful viewing prob.:  $\delta_n(S_n[t]) = 1$

$$W_1[1] = 0.5 * 0.5 + (0.2 * 3 + 1 * 0.3) * 1 = 1.15$$

$$W_2[1] = 0.5 * 0.5 + (0.2 * 2 + 1 * 0.6) * 1 = 1.25$$

$$W_3[1] = 0.5 * 0.5 + (0.2 * 1 + 1 * 1) * 1 = 1.45$$

Scheduled!

# Theoretical Bounds

$$\text{Algo: } S^*[t] \in \underset{S \in \mathcal{S}^{(c)}}{\text{argmax}} \sum_{n=1}^N (S_n[t]Q_n[t] + (\eta T_n[t] + K w_n) \delta_n(S_n[t]))$$

- Our proposed algorithm asymptotically optimizes the application-level throughput and provides seamless user experience guarantees while meeting the minimum service rate requirement, i.e.,

$$\lim_{L \rightarrow \infty} \frac{1}{L} \sum_{t=0}^{L-1} \sum_{n=1}^N E[w_n \delta_n(S_n[t])] \geq U^* - \frac{B(\eta)}{K} \quad ; \quad \lim_{L \rightarrow \infty} \frac{1}{L} \sum_{t=0}^{L-1} \sum_{n=1}^N U_n^* E[T_n[t]] \leq \frac{B(\eta) + KNw_{\max}}{\eta}$$

where  $B(\eta) \triangleq \sum_{n=1}^N \frac{r_n^2 + R_M^2}{2} + \eta N$ ,  $U^*$  is the optimal value of the optimization problem

- $K \nearrow$ , application-level throughput  $\nearrow$ , mean TSLS  $\nearrow$  (seamless user experience  $\searrow$ )
- $\eta \nearrow$ , application-level throughput  $\searrow$ , mean TSLS  $\searrow$  (seamless user experience  $\nearrow$ )

# Simulation

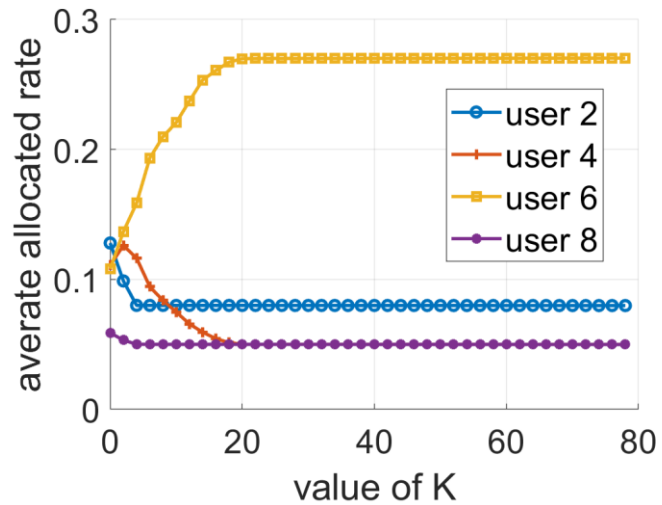
- 8 users
- Synthetic head motion data generated from the dataset [Bao, Wu, Zhang, Ramli, Liu, 2016]
- ON-OFF channel fading
- At most two users can be scheduled
- Total available rate: 1
- Rate set:  $\{0, 0.3, 0.4, 0.5, 0.7, 1\}$

	User 1	User 2	User 3	User 4
Required rate $r_n$	0.1	0.08	0.11	0.05
Weight $w_n$	0.2	0.1	1.0	0.8
Fading prob. $p_n$	0.8	0.9	0.7	0.9
	User 5	User 6	User 7	User 8
Required rate $r_n$	0.18	0.06	0.16	0.05
Weight $w_n$	0.9	1.2	0.3	0.2
Fading prob. $p_n$	0.8	0.9	0.7	0.8

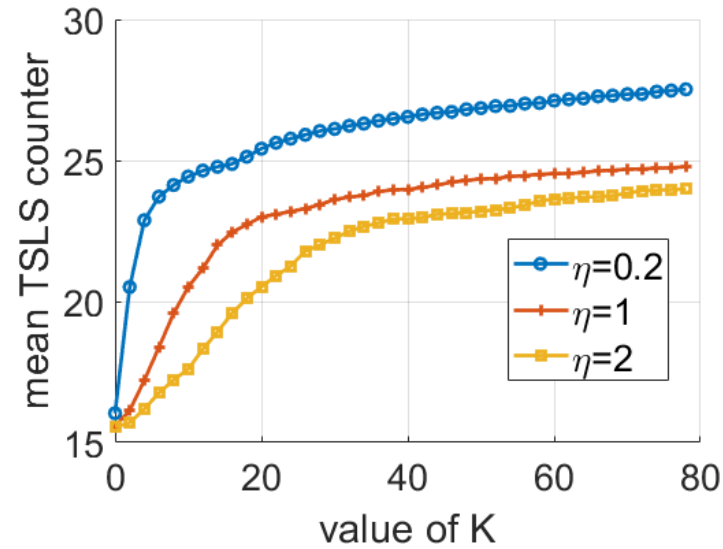
# Simulation (Cont')

Required rate:

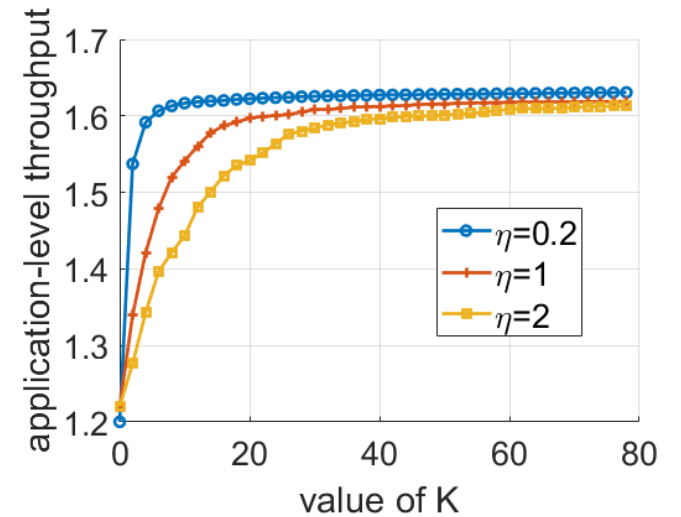
User	$r_n$
User 2	0.08
User 4	0.05
User 6	0.06
User 8	0.05



Average rate for each user:  $\eta = 1$



Average TSLS



Application-level throughput



# Related Work

- Panoramic Video Streaming:
  - [Bao, Wu, Zhang, Ramli, Liu, 2016]
  - [Qian, Han, Xiao, Gopalakrishnan, 2018]
  - [Perfecto, Elbamby, Del Ser, Bennis, 2020]
  - [Chen, Li, Srikant, 2020]
  - ...
- Wireless Scheduling Design:
  - [Li, Li, Eryilmaz, 2014]
  - [Neely, 2010]
  - [Hou, Kumar, 2013]
  - ...

# Conclusions

- The successful viewing probability as the function of the delivered portion
- A motion-prediction-based scheduling algorithm by integrating it into the stochastic network optimization framework
- The proposed algorithm can provide desired application-level throughput and service regularity guarantees
- Simulation results with real datasets demonstrated the efficiency of our proposed algorithm

Thank you!